Avoid Delay of Packets and Improve the Behavior of Switch/Router under Self-Similar-Type Variable Input Traffic

Anand Dohare¹, B.Mallikarjuna², Tulika³, D. Arun Kumar Reddy⁴

Mohd.Shahjad⁵, D. Mallikarjuna Reddy⁶

¹Research Scholar, SHUATS University, Allahabad, India-211007

²School of Computing Science & Engineering, Galgotias University, Greater Noida, Uttar Pradesh, India-

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³Assistant Professor, SHUATS University, Allahabad, India-211007
 ⁴Research Scholar, Bharathiar University, Coimbatore, Tamil Nadu, India-641 046
 ⁵Research Scholar, Uttarakhand Technical University, Dehradun, 248007
 ⁶Assistant Professor, GITAM University, Hyderabad, India-502 329

Abstract — This paper analyzes the switch/router under self-similar variable traffic with delay behavior on switch/router and the improvement in the performance of packet traffic under selfsimilarity nature of long-range dependency (LRD). The LRD property degrades the router performance. Consequently, it is essential for buffer conception of a switch in communication engineering technology. In this study, by taking into account voids, we examine the switch under self-similar input traffic. We simulated this model using queuing model and the MMPP (Markov-Modulated Poisson Process) and observed the variance in MMPP under selfsimilar traffic. We investigated traffic metrics, fitting specification and mean delay, which demonstrated improvement in the performance of switch/router.

Keywords — *MMPP*, *self-similar traffic*, *LRD*, *voids*, *mean delay*, *traffic intensity*.

I. INTRODUCTION

It observed that packet traffic principles of longrange dependent LAN and WAN are self-similar [1]. The qualities of service (QOS) parameters are numerous; however, we focused on packet delay and packet loss probability on routers and switches [2]. The Markovian Arrival Process (MAP) has been implemented in self-similar traffic [3]. MMPP is useful for calculating the probabilities at the time of packet transmission; hence, MMPP is a particular case of MAP, and interrupted poison process (IPP) has been implemented as a special case of MMPP [4]. The associate perception of MAP is Markov Modulated Poisson Process (-MMPP), where is superposition, having 2-state IPP. -MMPP contains the second-order statistics of MMPP and the selfsimilar traffic over desired time scales [5].

The switch and router have packet loss probability and mean packet delay with input or output traffic, the two main QoS parameters. L.P. Raj Kumar and K. Sampath Kumar (2010) designed queuing technique. it follows the principle of hurst index and traffic intensity [6]. It is examined by both analytical and simulation results. Packet lengths are assumed to follow exponential distribution [7]. But when the packet length varies, voids occur in the switch buffer and the performance of the switch degrades because voids incur excess load [8]. This issue is not addressed in other studies [6, 7]. In the study[8], packet arrivals are assumed to follow Poisson distribution which is proved to be unrealistic [9]. Here we are concerned with the performance of switch by applying the queuing technique, taking voids into account.

The rest of the paper describes the overview of self-similar processes, section 2 describes the related work on MMPP and then queuing method. Section 3 is about modelling of switch for variable length packets in buffers. Section 4 is on effective results being carried out by the queuing system, the packet loss probability and means waiting time and demonstrates the effects of parameters such as Hurst index, traffic intensity, number of components in superposition, and time scales. Section 4 concludes.

II. RELATED WORK

Krishna P.V, Misra S, Oaidat MS, Saritha V, N. Ch. S. N Iyengar (2009) proposed a sequencing technique on 802.11medium access control that improves the performance of wireless networks [19]. This method consists of two techniques: Request To Send (RTS) or Clear To Send (CTS). RTS/CTS method avoids wastage of the bandwidth, which is achieved by sequencing [18]. They proved this technique using NS2 toolkit. In mobile computing, QoS is required for channel allocation; and effective communication needs frequency spectrum that provides effective bandwidth spectrum [17]. They proposed QoS parameter hand off calls and drop off calls, maintaining queues for all types of calls [16]. They could not focus on communication devices such as switch and router. There is a necessity to improve the effective communication strategy on switch/router [15]. The network traffic on switch/router contains delay of packets and some of the malicious packets are dropped. It requires apply effective strategy for legitimate traffic, as shown in figure 1. and Γ are given in equation 2. The MMPP is fully parameterized by the infinitesimal generator Q of the Markov process and the vector



Figure 1: Model for switch or router in the network

Self-Similar Traffic - Markov Modulated Poisson Process (SSTMMPP) follows second-order selfsimilar process.

Algorithm 1: SSTMMPP Algorithm

Step 1: Let Y is a second-order stationary process with variance σ^2 , it divides the different time intervals ^[14]

$$Y = \{Y_t \mid t = 1, 2, 3, \dots, \}, \text{ if the packet arrived at } t^{th}$$

interval, then $Y^{(n)} = \{Y^{(n)}\}, \text{ where }$

 $Y_t^{(n)} = \frac{1}{n} \sum_{i=1}^n Y_{(t-1)n+i}, \quad t = -1, 2, 3...., \text{ m is the}$

average original sequence for overlapping packets. Then Y is the exact second-order

self-similar process with the Hurst parameter, $\Gamma = 1 - \delta / 2$, if

$$Var(Y^{(n)}) = \sigma^2 n^{-\beta}, \quad \forall n \ge 1.$$

Step 2: There is doubly stochastic process as per Morkovian Modulated Poison. This process in which the arrival rate is given by $\chi[J_{i}]$,

here $J_i, t \ge 0$ is an n-state Markov process^[13]. The arrival rate can therefore take on only n values, namely, $\chi_1, \chi_2, \dots, \chi_n$. It is equal to χ_j whenever the Markov process is in the

state J, $1 \le j \le n$.

Let Assume Γ be the diagonal matrix with $\Gamma_{jj} = \chi_j, \ 1 \le j \le m$. In the case of two states, Q

$$\chi = (\chi_1, \chi_2, \dots, \chi_m)$$
 of the arrival rates ^[12]

$$Q = \begin{bmatrix} -S_1 & S_1 \\ S_2 & -S_2 \end{bmatrix} \quad \Gamma = \begin{bmatrix} \chi_1 & 0 \\ 0 & \chi_2 \end{bmatrix}.$$
(2),

Step 3: χ is the mean arrival rate of MMPP, which \vec{x}

is given by $\chi = \vec{\pi} \Gamma e$, here

Q is the stationary probability vector contains $\vec{\pi}$, where $\vec{\pi}Q = 0, \vec{\pi}e = 1$ and e.

It contains -1 of designated dimension column vector. If we let N_t , t ≥ 0 , in the

interval (0, t] contains the number of arrivals, for the stationary MMPP, the

mean of N_t is:

$$\frac{S_2 \chi_1 + S_1 \chi_2}{S_1 + S_2} tmean = \mu = E[N_t].$$
 (3)

The variance of N_t is:

$$E[N_t] + \frac{2S_1S_2(\chi_1 - \chi_2)^2}{(S_1 + S_2)^3}t - \frac{2S_1S_2(\chi_1 - \chi_2)^2}{(S_1 + S_2)^4}[1 - e^{-(S_1 + S_2)t}] = Var[N_t] = \sigma^2.$$
(4)

III. SELF-SIMILAR TRAFFIC MMPP/M/1/K QUEUING METHOD

In queuing system, asynchronous switch followed the self-similar structure of variable length packet input traffic is modeled. In system, the packets are allowed into the switch MMPP of states which can be written as matrix representation , where are matrices[11]. The self-similar traffic Markovian-modulated Poisson process (MMPP) MMPP/M/1/K following algorithm.

Algorithm 2: MMPP/M/1/K Algorithm

Step 1: μ is the rate of service exponential value. It defines $n \times n$ which is the order of matrix E_l , $l \ge 0$ contains (i, j) elements, the probability of packets at the time of departure is 0. At the initial stage, one departure state is i, the next departure state is j and k is a throughput of the arrivals. E_k is defined as follows [10].

$$\sum_{l=0}^{\infty} E_{k} z^{k} = \mu \int_{0}^{\infty} e^{lQ - \Gamma + \Gamma z J x} e^{-\mu x} dx \quad (5)$$
$$\sum_{l=0}^{\infty} E_{l} z^{l} = \mu \left(\mu I - (Q - \Gamma + \Gamma z) \right)^{-1}$$
$$\sum_{l=0}^{\infty} D_{l} z^{l} = \mu^{2} \sum_{l=0}^{\infty} \left(\frac{Q - \Gamma + \Gamma z}{\mu} \right)^{l}, \quad (6)$$

Step 2: Let I be the unit matrix chosen by the dimension. Equation 6 extends the distribution exponential service, where in equation (6) can be written as:

$$E_0 = \mu^2 \left(I + \sum_{r=0}^{\infty} \left(\frac{Q - \Gamma}{\mu} \right)^r \right).$$
(7)

For $l, n \ge 1$, let T(n, l) be the coefficient of z^{l-1} in $\mu^2 \left(\frac{Q-\Gamma+\Gamma z}{\mu}\right)^n$, n^{th} from the equations (6) and (7), which can be written as:

$$T(1,1) = \mu^2 \left(\frac{Q-\Gamma}{\mu}\right), \quad T(1,2) = \mu^2 \frac{\Lambda}{\mu}$$
$$T(n,1) = \mu^2 \left(\frac{Q-\Gamma}{\mu}\right)^n, \quad T(n,l) = 0, \text{ if } l > n+1,$$
And

 $T(n,1) + T(n,2)z + T(n,3)z^{2} + \dots + T(n,l)z^{l-1} + \dots + T(n,n+1) = \mu^{2} \left(\frac{Q - \Gamma + Az}{\mu}\right)^{2}$

 $\mu^2 \left(\frac{Q-\Gamma+\Gamma_z}{\mu}\right)$, it obtains the series of equation as

follows:

 $[T(n,1) + T(n,2)z + T(n,3)z^{2} + \dots + T(n,l)z^{l-1} + \dots + T(n,n+1)z^{n}]\mu^{2}\left(\frac{Q-\Lambda+\Lambda z}{\mu}\right)$ $= T(n+1,1) + T(n+1,2)z + \dots$

Associate the coefficients of power of Z and then

$$T(n+1,1) = T(n,1)\mu^{2} \left(\frac{Q-\Gamma}{\mu}\right),$$

$$T(n+1,2) = T(n,2)\mu^{2} \left(\frac{Q-\Gamma}{\mu}\right) + T(n,1)\mu^{2} \frac{\Gamma}{\mu},$$

$$T(n+1,q) = T(n,q)\mu^{2} \left(\frac{Q-\Gamma}{\mu}\right) + T(n,q-1)\mu^{2} \frac{\Gamma}{\mu}$$

$$q \in \mathbb{N}.$$
 Equation (9) is formed to date the equations

is formed to date the equations

(6), (7), and (8)

$$E_{s} = \sum_{k=s}^{\infty} T(k, s+1), s = 1, 2, \dots$$
(9)

Step 3: The recurrent relation matrix R from equations 7, 8, and 9 have recurrent relations lead to obtaining equation 10. Then we framed Markov chain of each packet that arrived in the switch or router $\{L(n), \tilde{J}(n) / n \ge 0\}$ at the leaving epochs of the packets departure rate in the queue arrangement MMPP / M/1/k, and J(n) denotes the situation is MMPP, the packets arrived in the entire state space

 $S = \{(b,i)/0 \le b \le K - 1, 1 \le i \le m\},\$ where

L(n) denotes buffer tenancy:

$$R = \begin{bmatrix} HE_0 & HE_1 & \dots & HE_{K-2} & HF_{K-1} \\ E_0 & E_1 & \dots & E_{K-2} & F_{K-1} \\ 0 & E_0 & \dots & E_{K-3} & F_{K-2} \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & E_1 & F_2 \\ 0 & 0 & \dots & E_0 & F_1 \end{bmatrix}$$
(10)

Equation (10) matrix R provides each packet pertaining under self-similar variable embedded Markov chain that transition to probability matrix

where $H = (\Gamma - Q)^{-1}\Gamma$, consisting of conditional probabilities that system is not busy and

$$F_i = \sum_{l=i}^{\infty} E_l$$
. Let $\overline{y_k}$, $(0 \le k \le K - 1)$

be an $1 \times n$ vector whose i^{th} element is the number of packets in the system, k is a stationary conditional probability and i is the packet that arrived at random. At the end, the packet loss probability (PLP) has been defined [11] is

$$PLP = 1 - \frac{\left(1 - \overline{y_0}e\right)}{\rho} \qquad (11)$$

where $\rho = \frac{\lambda}{u}$, traffic intensity and λ is the mean arrival rate of MMPP where

 $\lambda = \pi \Lambda e$

A. Modeling of Switch for Variable Length Packets in Buffers

In this segment, we consider the switch buffer for uneven extent packets. In this buffering arrangement, the second packet arrives at the time when the first one is transmitted. In order to avoid disputation, the second packet is to be deferred by physical (8) mechanism, namely, fiber delay lines (FDL). Let the first packet arrive at the system at time t, tf is the time at which the packet is transmitted and this

packet will be delayed by an amount of time tf -t. In the ocular buffer system, the delay time must be a consecutive multiple of delay line unit (DLU) denoted by D as shown in the Fig. 1. Hence, the actual delay time, (say) becomes larger than tf - t and is given by.

$$\Delta = \left[\frac{t_f - t}{D}\right] D, \qquad (12)$$

where [x] represents the smallest integer not greater than x. Therefore, $\tau = \Delta - t_f + t \ge 0$ is the time interval when the system is idle. The waiting time τ is called void which automatically increases when the packet length increases; hence, it degrades the system performance. Several methods have been proposed to resolve this problem. The possible method is the void filling scheme $^{[16, 20, 21]}$. In this context, appropriate models using both infinite and finite queuing systems are proposed by Callegati^[8]. In the said models, input traffic is assumed to follow Poisson process which is not realistic, because Poisson process cannot emulate the self-similar nature of the traffic. Instead, we have an MMPP/M/1/k queuing system to analyze the switch.

We assume that actual packet length \mathcal{V} follows exponential distribution with parameter \mathcal{H} and $\overline{\mathcal{V}}$ its average. Let \mathcal{V}_e be the excess length and $\overline{\mathcal{V}_e}$ be its average. We assume that void length τ follows rectangular distribution and takes the values between 0 and \mathcal{D} . Therefore, its average is $\overline{\tau} = \frac{\mathcal{D}}{2}$.



Figure 2: Illustration of Queuing in a FDL buffers

we derive equation (13) from equation (12).

 $\overline{v_e} = \overline{v} + (1 - PLP)\pi_q \frac{D}{2}$ (13) where π_q is the probability that a packet not lost is queued. The traffic intensity in this case would be $\rho_e = \lambda \overline{v_e}$ and this ρ_e is called excess load. We know that if X is exponential and Y is uniform, then the density of Z = X + Y is given by

$$h(z) = \begin{cases} 0 & z \le 0 & (14) \\ 1 - e^{-z} & 0 \le z \le 1 \\ (e - 1)e^{-2} & z \ge 1 \end{cases}$$

Hence, the excess length V_e is not exponentially distributed, but it has the density function given above that would greatly increase the complexity in (5). For this reason, we assume that packet length distribution follows exponential with parameter $\mu_e = \frac{1}{V_e}$, this simplifies the problem and we have $\rho_e = \frac{\lambda}{\mu_e}$.

In the case of *MMPP* / *M* / 1 / *k* queuing system, we have $\pi_q = 1 - (1 - \rho) = \rho$

As discussed in [7], starting with ρ , the packet loss probability PLP can be calculated using eqn (11). Using this PLP, $\overline{\nu_e}$ can be computed using eqn. (13). Finally, the mean delay (*w*) is given by

$$w = \overline{v}_e \sum_{k=1}^{\infty} k(\overline{y}_k \ e). \quad (15)$$

IV. RESULTS AND DISCUSSIONS

The results describe the intensity of traffic, time scale of packet transmission, Hurst index parameter . Primarily the transition rate of matrix Q and packets arrival rate in MAP are represented as Λ . S.K. Shao, Malla Reddy Perati (2005) gave MMPP design model in the form of generalized variancebased method for the self-similar traffic. It contains H=0.7,0.8,0.9, Hurst parameter values variance $\sigma^2 = 0.6$, arrival rate $\lambda = 1$ with different time intervals $[10^2, 10^5], [10^2, 10^6],$ and $[10^2, 10^7]$. The stationary probability vector ^y is enumerated and then mean delay w is computed against traffic intensity without taking voids in ^[7, 22]

$$w = \overline{v} \sum_{k=1}^{\infty} k(\overline{y}_k e)$$
 to r

using the formula $\overline{k=1}$ to realize the effect of void consideration. From Fig.2, it is apparent that as traffic intensity increases, the mean delay increases. First, at $\rho = 0.4$, the packet loss probability is computed ^[12], then $\overline{V_e}$ is computed using eqn (13), and then the mean delay is computed against D using eqn (15). Results are depicted in Fig. 3.

Fig.4 illustrates that mean delay automatically increases D and the values are higher when H and time scale are higher. Fig. 5 depicts the variation in mean delay against traffic

intensity ρ , when the Delay Line Unit DLU (D)=0.1. From Fig. 4, one can infer that mean delay increases as ρ increases. From Figs.4 and 5, it is clear that the trend in the case of voids is different from the one without voids, and the values are higher in the case of voids..



Figure 3: Mean Waiting Time with Traffic Intensity when H=0.7 and d=4.



Figure 4: Variation in Mean Waiting Time with Delay Line Unit at different Γ .



Intensity at various Γ .

V. CONCLUSIONS

The switch/router contains uneven length packet input traffic, which can be modeled as self-similar MMPP / M / 1 / k method, and we analyzed delay behavior of the switch. In this proposed paper of MMPP model, the variance in self-similar traffic over a time-scale has to be matched. In this model, packet lengths and voids are assumed to follow exponential distribution and rectangular distribution. Our numerical results reveal that time scale and Hurst parameters do have an impact on the mean delay. If the factors are Γ , ρ and DLU(D), the mean delay increases. Finally we understand that one could select the appropriate time interval D to approach the QoS obligation. This type of analysis is supportive evidences to switch/router. It is dimensioning not only the switch/router but also Hub and any network device. The results are based on the concept of self-similar variable length packet traffic. In future work to implement Markovian models in cloud computing, big data and internet of things [23] [24] [25] [26].

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